

Counting orbits on colourings and flows

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- ▶ Real chromatic roots are dense in $[\frac{32}{27}, \infty)$ (Thomassen)
- ▶ Complex chromatic roots are dense in \mathbb{C} (Sokal)

Orbit-counting Lemma

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Theorem

Let G act on a set X . Then the number of orbits of G on X is equal to the average number of fixed points on X of the elements of G :

$$\#\text{Orbits}(G, X) = \frac{1}{|G|} \sum_{g \in G} \text{fix}_X(g).$$

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Said otherwise, the number of orbits is the expected number of fixed points of a random element of G .

Orbital chromatic polynomial

Let g be an automorphism of a graph Γ . Denote by Γ/g the graph obtained by shrinking every cycle of g to a single vertex. The number of k -colourings of Γ fixed by g is equal to the number of colourings of Γ/g . For a colouring is fixed by g if and only if it is constant on the cycles of g (and so induces a proper colouring of Γ/g).

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So, if G is a group of automorphisms of Γ , define the **orbital chromatic polynomial** of Γ and G to be

$$OP_{\Gamma,G}(x) = \frac{1}{|G|} \sum_{g \in G} P_{\Gamma/g}(x).$$

Then for positive integers k , the number of orbits of G on the k -colourings of Γ is $OP_{\Gamma,G}(k)$.

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Here is an example to show that, unlike chromatic roots, orbital chromatic roots can be negative.

Take Γ to be the null graph on n vertices and G the symmetric group S_n . Then an orbit on k -colourings is a choice of n things from a set of k , where repetitions are allowed and order is not significant. This number is $\binom{n+k-1}{n}$. So

$$OP_{\Gamma,G}(x) = \frac{1}{n!} x(x+1) \cdots (x+n-1),$$

with roots $0, -1, \dots, -(n-1)$.

Orbital chromatic roots, continued

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Proof.

Take Γ to consist of m triangles and G to be the symmetric group S_m . Then

$$OP_{\Gamma,G}(x) = \frac{1}{m!} q(x)(q(x) + 1) \cdots (q(x) + m - 1),$$

where $q(x) = x(x - 1)(x - 2)$.

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The equation $x(x - 1)(x - 2) = -k$ has a unique negative root α_k , and the spacing of the α_k becomes denser as k increases.

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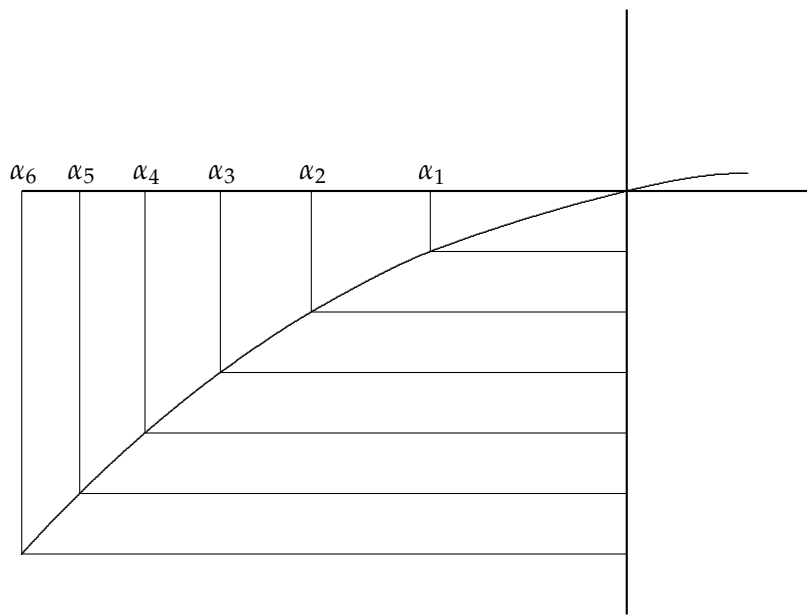
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The equation $x(x - 1)(x - 2) = -k$ has a unique negative root α_k , and the spacing of the α_k becomes denser as k increases.

Now by taking the join with a complete graph of size s , we can translate the roots to the right by any integer s . □

Orbital chromatic roots, continued



Parity, 1

Theorem

The sign of $P_\Gamma(x)$ is $(-1)^v$ for $x < 0$, $(-1)^{v+c}$ for $x \in (0, 1)$, and $(-1)^{v+c+b}$ for $x \in (1, \frac{32}{27}]$, where v, c, b denote the numbers of vertices, connected components, and blocks of Γ .

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Since a permutation is even if and only if the numbers of points and cycles are congruent mod 2, we have:

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Theorem

Let V and C be the sets of vertices and connected components of Γ .

- ▶ *(a) Suppose that every element of G is an even permutation of V . Then $OP_{\Gamma,G}(x)$ has no roots in $(-\infty, 0)$.*
- ▶ *(b) Suppose that every element of G is an even permutation of $V \cup C$. Then $OP_{\Gamma,G}(x)$ has no roots in $(0, 1)$.*

Parity, 2

The analogous result for blocks is false. The reason is that, while vertices and connected components of Γ/g correspond to cycles of g on vertices and connected components of Γ , the analogous statement is not true for blocks.

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Under these hypotheses, the only root less than 1 is 0.

Flow and tension polynomials

Take a fixed but arbitrary orientation of the edges of Γ . A **flow** on Γ with values in the abelian group A is a function from the oriented edges of Γ to A with the property that the net flow into any vertex is zero (calculated in A).

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The flow and tension polynomials are specialisations of the **Tutte polynomial** of Γ .

Orbital flow polynomial

Things are more complicated for orbits on flows: the answer does depend on the structure of A . Precisely, given a group G of automorphisms of the graph Γ , there is a polynomial $OF_{\Gamma,G}(x_0, x_1, x_2, \dots)$ in indeterminates indexed by \mathbb{N} , such that the number of G -orbits on nowhere-zero A -flows on Γ is $OF_{\Gamma,G}(a_0, a_1, \dots)$, where a_i is the number of solutions of $ix = 0$ in the abelian group A . Note that $a_0 = |A|$ and $a_1 = 1$.

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This orbital flow polynomial is a specialisation of an **orbital Tutte polynomial** of Γ and G , which involves two potentially infinite families of variables.

If the variable x_i actually occurs, then G must contain an element of order i . Thus we recover Tutte's observation by taking G to be the trivial group.

Orbital flow roots

An **orbital flow root** is a root of the polynomial $OF_{\Gamma,G}^1$ obtained by putting $x_i = 1$ for all $i > 0$, for some pair (Γ, G) . It counts orbits on flows in the case where $\gcd(|A|, |G|) = 1$.

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The method of proof also shows that every value $1/k$, for $k \in \mathbb{N}$, is a limit point of real orbital flow roots.

The limitation in the method is that there is no way to “translate” orbital flow roots to the right, as there is for orbital chromatic roots!

Invariant factors and duality, 1

Let R be a principal ideal domain. Given an $m \times n$ matrix M over R , we define the **row space** of $\rho(M)$ and the **null space** $\nu(M)$ as usual:

$$\rho(M) = \{yM : y \in R^m\},$$

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M can be put into *Smith normal form* by elementary row and column operations: this is a matrix with r non-zero diagonal elements d_1, \dots, d_r and all other entries zero, where d_i divides d_{i+1} for $i = 1, \dots, r-1$. The elements d_1, \dots, d_r are uniquely determined up to multiplication by units of R . They are the *invariant factors* of M . By convention, we also take 0 to be an invariant factor with multiplicity $n - r$, so that there are n invariant factors in all.

Invariant factors and duality, 2

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- ▶ M has a dual;
- ▶ all invariant factors of M are zero or units;
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- ▶ M has a dual;
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- ▶ M is equivalent (by elementary row and column operations) to a totally unimodular matrix.

If Γ is a graph with oriented edges, and M and M^* are its signed vertex-edge and cycle-edge incidence matrices, then M and M^* are dual.

Orbital Tutte polynomial, 1

Assume that (M, M^*) is a dual pair over a principal ideal domain R . The linearly independent sets of columns of M are the independent sets of a **matroid**. The linearly independent sets of columns of M^* form the **dual matroid**.

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If g is an automorphism of M (represented as an $n \times n$ matrix), and 1 is the identity matrix, set

$$M_g = \begin{pmatrix} M \\ g - 1 \end{pmatrix}, \quad M_g^* = \begin{pmatrix} M^* \\ g - 1 \end{pmatrix}.$$

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For any subset S of $E = \{1, \dots, n\}$, and any matrix N with n columns, we let $N[S]$ be the submatrix of N consisting of the columns with indices in S .

Orbital Tutte polynomial, 2

Take two sets $(x_i : i \in I)$ and $(x_i^* : i \in I)$ of indeterminates, where the index set I is the set of associate classes in R . For any matrix N , let $x(N)$ be the monomial defined as follows: take the invariant factors of N (completed with zeros so that the number of them is equal to the number of columns of N), and multiply the corresponding indeterminates. Define $x^*(N)$ similarly, using the other set of indeterminates.

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Now let G be a finite group of automorphisms of M , and define the **orbital Tutte polynomial** $OT(M, G)$ in the indeterminates $(x_i, x_i^* : i \in I)$ as follows:

$$OT(M, G) = \frac{1}{|G|} \sum_{g \in G} \sum_{S \subseteq E} x(M_g[S]) x^*(M_g^*[E \setminus S]).$$

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Theorem

If G is the trivial group, then $OT(M, G)$ involves only x_0, x_1, x_0^* and x_1^* ; the substitution $x_1 = x_1^* = 1, x_0 = y - 1, x_0^* = x - 1$ gives the Tutte polynomial of M .

Graphs

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Theorem

Let M be the incidence matrix of a graph Γ over \mathbb{Z} , and let G be a group of automorphisms of Γ . Let A be a finite Abelian group. Then the substitution $x_i \leftarrow \alpha_i(A)$, $x_i^ \leftarrow -1$ (for all i) in $OT(\Gamma, G)$ gives the number of G -orbits on nowhere-zero A -flows on Γ , while the substitution $x_i \leftarrow -1$, $x_i^* \leftarrow \alpha_i(A)$ gives the number of G -orbits on nowhere-zero A -tensions on Γ .*

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Theorem

Let Γ be a connected graph. Then the orbital chromatic polynomial of $(\Gamma, G; k)$ is obtained from the orbital tension polynomial by substituting $x_0^ = k$, $x_i^* = 1$ for $i > 0$, and multiplying by k .*

Supports

Theorem

- ▶ Let f_m be the number of G -orbits on A -flows supported on precisely $n - m$ edges of Γ . Then

$$\sum f_m x^m = OT(\Gamma, G; x_i \leftarrow \alpha_i(A), x_i^* \leftarrow x - 1).$$

- ▶ Let t_m be the number of G -orbits on A -tensions supported on precisely $n - m$ edges of Γ . Then

$$\sum t_m x^m = OT(\Gamma, G; x_i \leftarrow x - 1, x_i^* \leftarrow \alpha_i(A)).$$

Orbital weight enumerator

A linear code over $\text{GF}(q)$ is the row space of a **generator matrix** M over $\text{GF}(q)$, and is the null space of a **parity check matrix** M^* (the generator matrix of the dual code). These matrices are duals; so if G is a group of automorphisms of C , the orbital Tutte polynomial P is defined as before. Since $\text{GF}(q)$ is a field, P involves only the variables x_0, x_1, x_0^*, x_1^* .

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The **orbital weight enumerator** is the homogeneous polynomial

$$W_{C,G}(X, Y) = \sum_{i=0}^n a_i X^{n-i} Y^i,$$

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Theorem

Let G be an automorphism group of a linear code over $\text{GF}(q)$. Then the orbital weight enumerator C is obtained from the orbital Tutte polynomial by the substitution

$$x_0 = x_1 = X - Y, \quad x_0^* = qY, \quad x_1^* = Y.$$

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